

Typing Dynamic Layer Composition

Atsushi Igarashi (Kyoto U.)

joint work with

Robert Hirschfeld (HPI)

Hidehiko Masuhara (Tokyo Tech.)

Hiroaki Inoue (Kyoto U.)

Context-Oriented Programming (COP)

Language [Costanza, Hirshfeld DLS05]

[Hirschfeld, Costanza, Nierstrasz JOT08]

Goal: Support for modularization of *behavioral variations* depending on the *dynamic context* of execution

Example: Mobile email app



When network is fast
inline images are shown



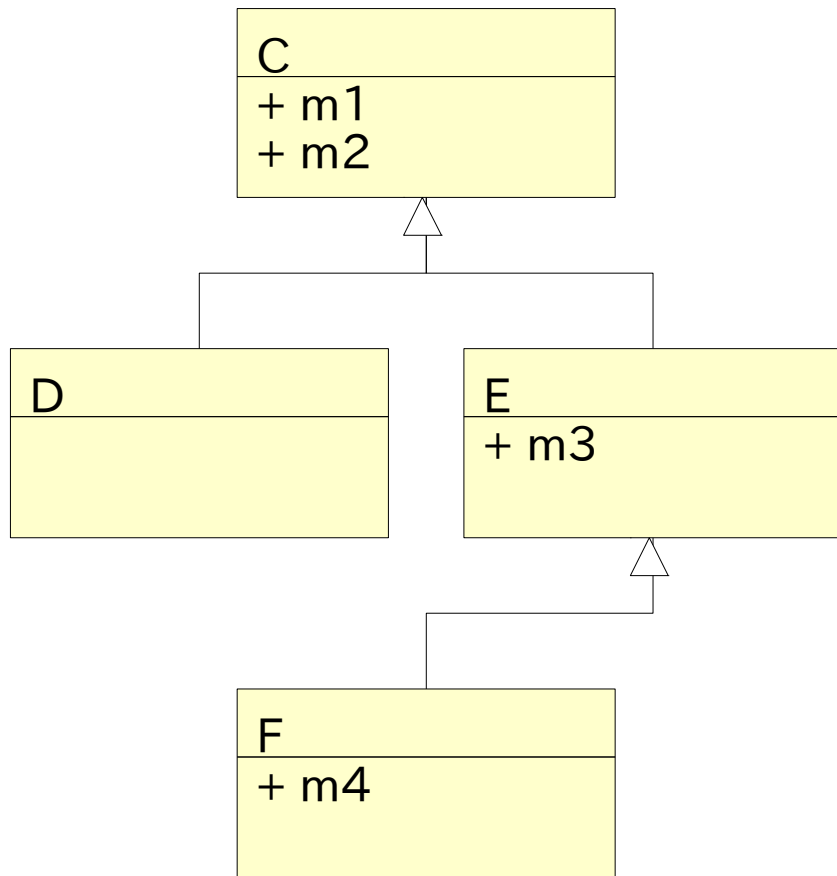
When network is slow
no images are shown

Common COP language features

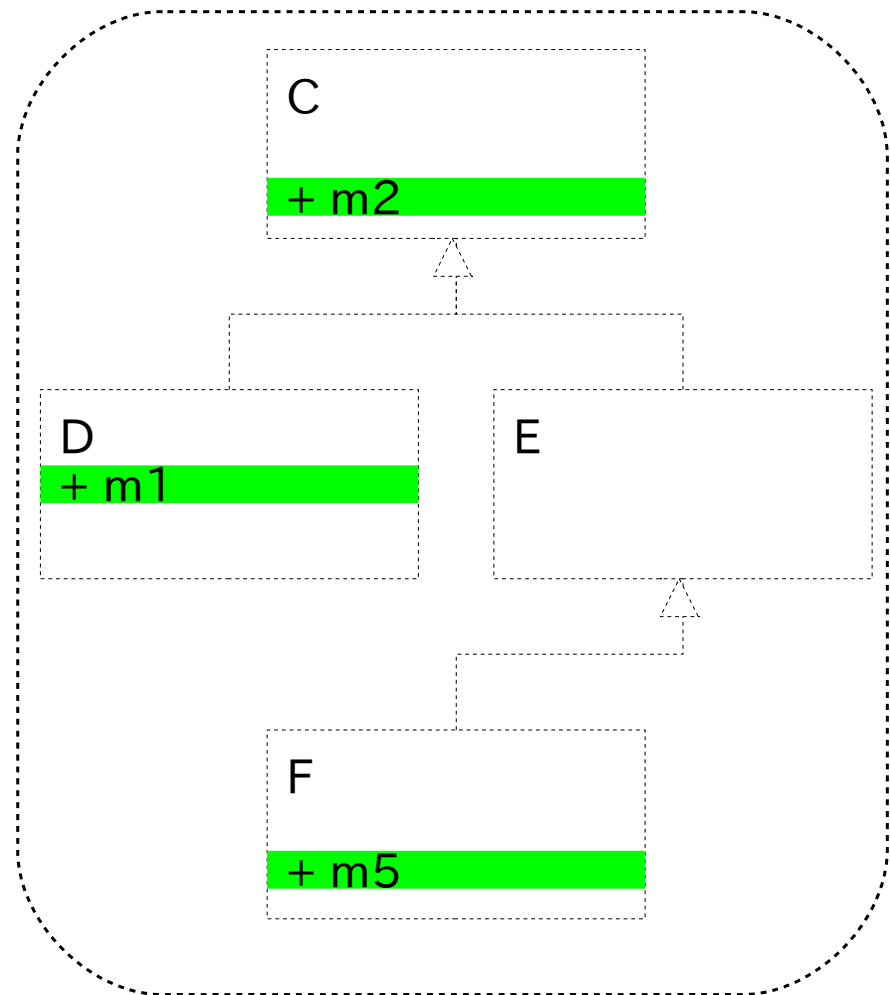
- Layer
 - A unit of behavioral variations, consisting of *partial* method definitions for multiple classes
 - (Loose) correspondence to contexts
 - A unit of cross-cutting modularity
- Dynamic layer activation
 - To change the behavior of a set of objects at the same time

Dynamic Layer Activation in COP

Base class hierarchy

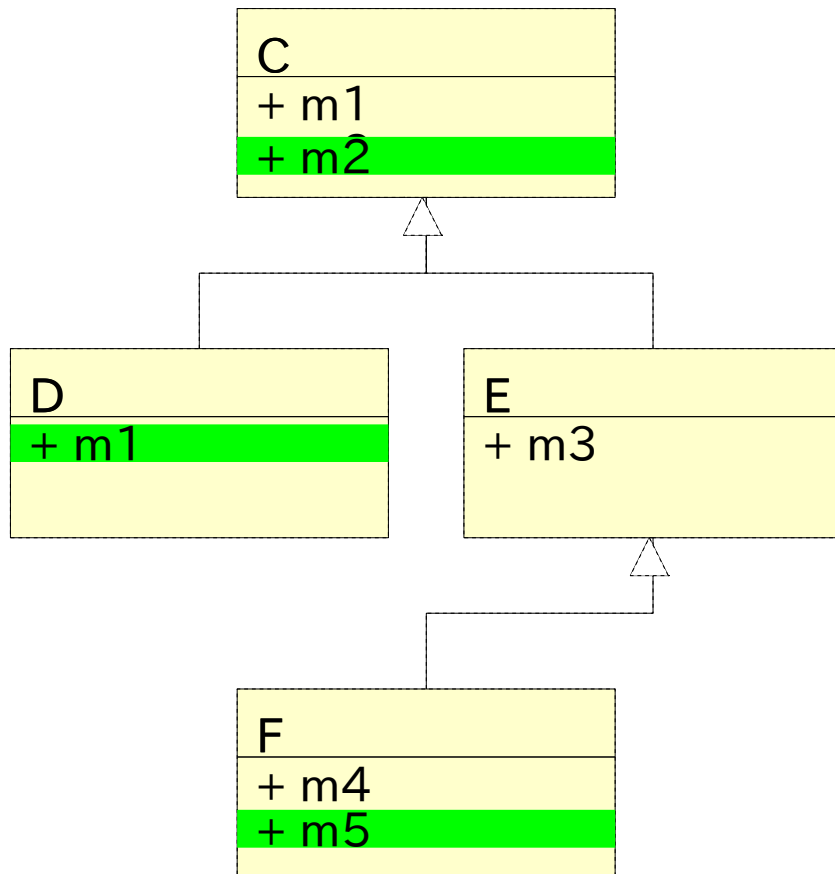


Layer of partial methods



Dynamic Layer Activation in COP

Base class hierarchy



- Layer activation changes behavior of objects *that have been already instantiated*
- Partial methods can call the original behavior by **proceed()**

This Talk

- Quick tour on JCop [Appeltaufer+], a specific implementation of COP on top of Java
 - With a more concrete example
 - (Comparison with AOP using pointcut/advice)
- Foundations of COPL

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Example: Telecom simulation

(adapted from AOP example)

- Class Conn to represent connection between two Customers
 - complete () when a connection has been established
 - drop () when the customers are disconnected
- Behavioral variations to consider
 - Recording the lengths of conversations
 - Billing

Base Program

```
class Conn { // Connection
    Conn(Customer a, Customer b) { ... }
    void complete() { ... }
    void drop() { ... }
    // details are not important ...
}
```

```
static Conn simulate() {
    Customer robert = ..., hidehiko = ...;
    Conn c = new Conn(robert, hidehiko);
                                // Robert calls Hidehiko
    c.complete(); // Hidehiko accepts
    c.drop(); // and hangs up
    return c;
}
```

Layer for Measuring Time

```
layer Timing {  
  Timer timer = ...;  
  void Conn.complete() { proceed(); timer.start(); }  
  void Conn.drop() { timer.stop(); proceed(); }  
  int Conn.getTime() { return timer.getTime(); }  
}
```

- The two methods in Conn are **modified** by *partial* method definitions to operate the timer
 - The original behavior is represented by `proceed()`
- `getTime()` is **newly introduced**
 - but also called “partial” method

Layer Activation with `with`

```
with (new Timing()) { // layer activation!  
    Conn c = simulate();  
    System.out.println(c.getTime());  
}
```

- `with` block to activate a layer
- Activation is effective even in methods invoked inside the block
- A layer *instance* has to be created
 - Layer instances are also first-class objects

Layer for Billing

```
layer Billing {  
    void Conn.drop() { proceed(); charge(); }  
    void Conn.charge() { ... getTime(); ... }  
}
```

```
with (new Timing()) {  
    with (new Billing()) {  
        Connection c = simulate();  
    }  
}
```

- Recently activated layer has priority
 - `drop()` will stop the timer, hang the call, and charge

Not in this example, but...

- One layer can contain partial methods belonging to different classes
 - c.f. Mixin layers [Smaragdakis&Batory 98]
- `super()` is also supported
- Layer inheritance/subtyping

Layer Inheritance/Subtyping

- Implementation of different billing policies, switched by run-time conditions

```
abstract layer AbsBilling {  
    void Conn.drop();  
    void Conn.charge();  
}
```

```
layer Billing1 extends AbsBilling { ... }
```

```
layer Billing2 extends AbsBilling { ... }
```

```
AbsBilling b =
```

```
    some_cond ? new Billing1() : new Billing2();
```

```
with(b) { ... }
```

Very rough Comparison with Aspect/J-style AOP

	COP	AOP
Unit of behavior	partial meth.	advice
Oblivious?	No	Yes
Join points	Meth. exec.	Many kinds
Pointcut	cflow + execution	Many kinds

Some Foundational Questions

- What is the semantics of method invocations?
 - What happens when the same layer is activated more than once?
 - How do `proceed`, `super`, and `with` interact with each other?
- How can types prevent `NoSuchMethodError`?
 - Object interface can change dynamically!
 - Only overriding partial methods can `proceed`

This Talk

- Quick tour on *COP* language features
 - With a more concrete example
- Foundations of *COPL*

A core calculus of COP: ContextFJ

[Hirschfeld, I., Masuhara FOAL'11]

ContextFJ = Featherweight Java [I.,Pierce,Wadler'99]

- + partial methods
- + `proceed()`, `super()`
- + with expressions
- layers are global and second-class
- no layer inheritance

ContextFJ<:

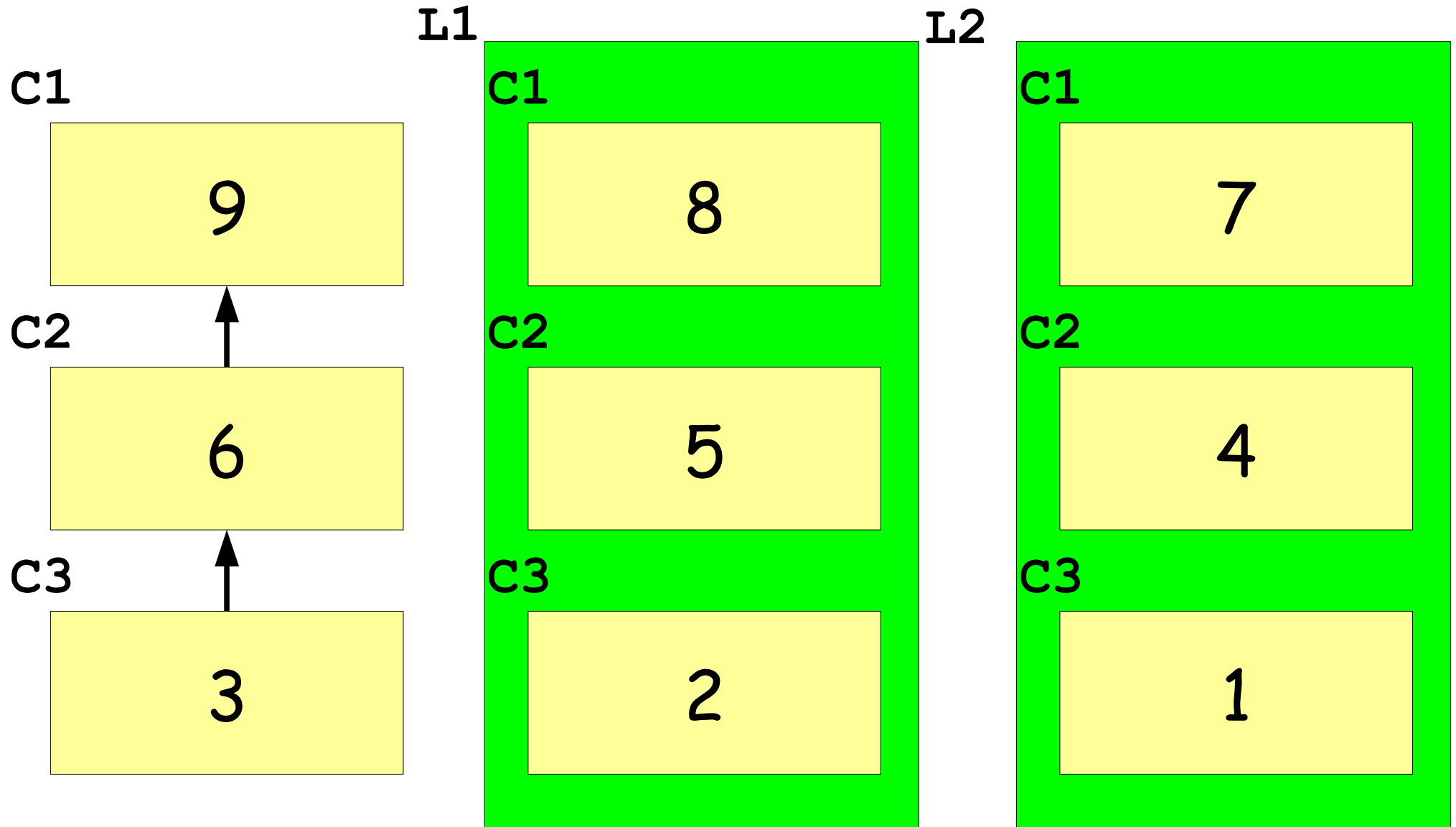
[Inoue&I. APLAS'15]

ContextFJ<: = Featherweight Java

- + partial methods
- + `proceed()`, `super()`
- + with expressions
- + first-class layers (w/o fields)
- + layer inheritance
- + layer subtyping

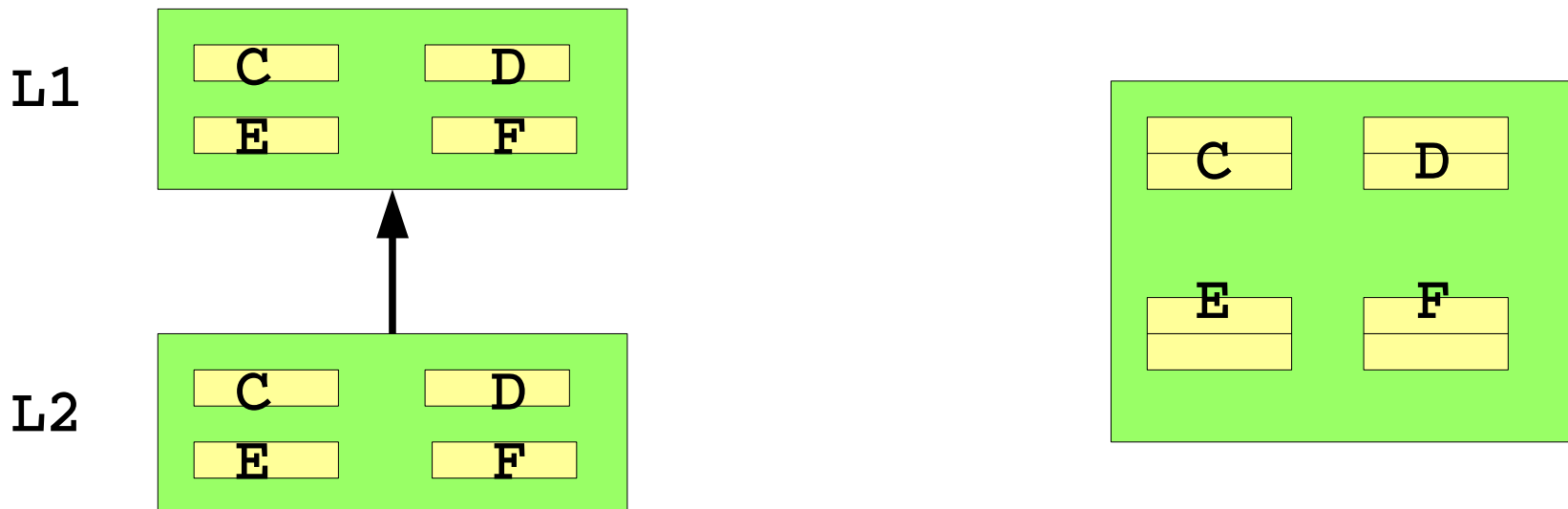
Semantics of Method Dispatch w/o Layer Inheritance

```
with (new L1()) {  
  with (new L2()) {  
    c3.m(...);  
  }  
}
```



Semantics with Layer Inheritance

- "3D" dispatching
- Each layer can be thought of as the result of (possibly overriding) composition of superlayers



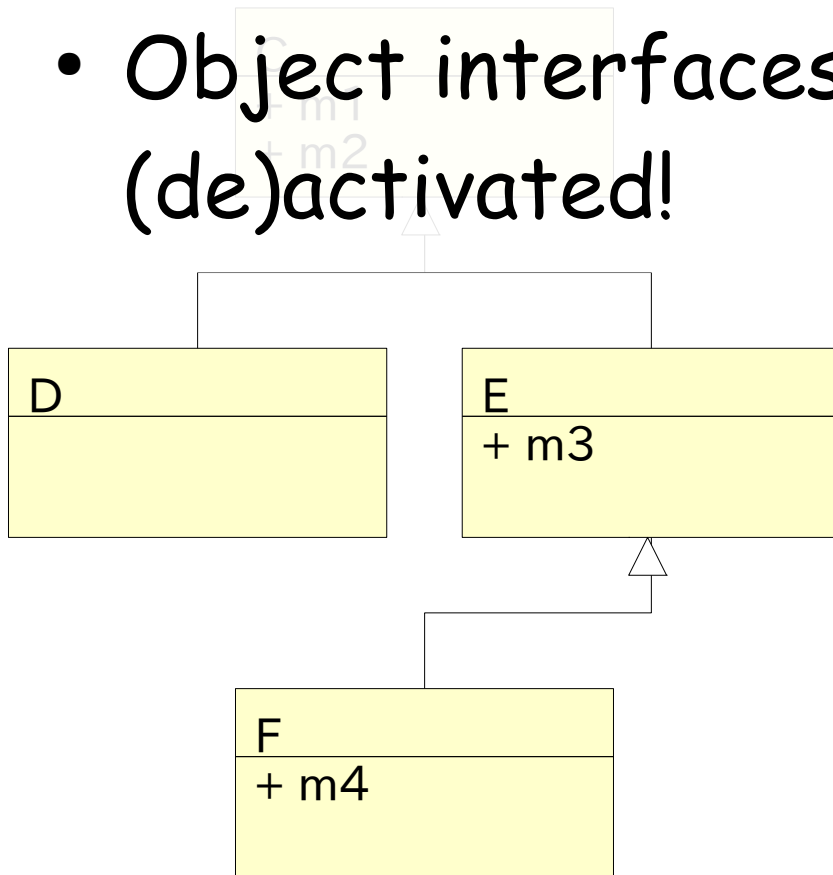
This Talk

- Quick tour on COP language features
 - With a more concrete example
- Foundations of COPL
 - (Operational) Semantics
 - Type System
 - To prevent “NoSuchMethodError” including dangling proceed calls

"Sounds like an old problem.

What is a challenge?"

- Object interfaces can change as layers are (de)activated!



Overriding partial method

+ m1

"Baseless" partial method,
which can dynamically change
the object interface!

+ m5

Key Ideas (1/2)

Approximating activated layers at each program point

- With the help of explicit “requires” declarations to specify inter-layer dependency
 - (Static analysis could dispense with such explicit declarations)

Key Ideas (2/2)

Two kinds of substitutability for layers

- When one layer L1 requires layer L2, does a sublayer of L2 can satisfy L1's requirement?
- When is it safe to pass an instance of a layer to where a supertype is expected?

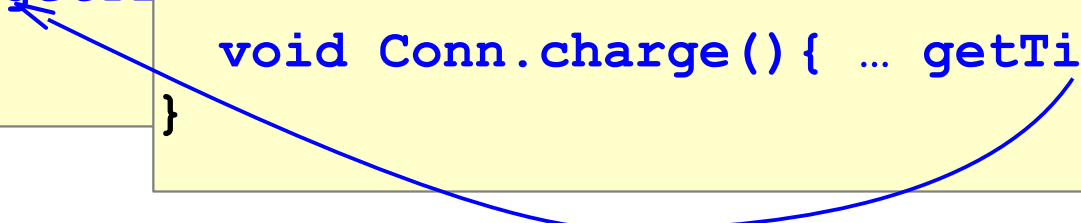
should be distinguished

Telecom example, revisited

```
class Conn {  
  Conn(Custome  
  void complet  
  void drop()  
}
```

```
layer Timing {  
  Timer Conn.tim  
  void Conn.comp  
  void Conn.drop  
  int Conn.getTi  
}
```

```
layer Billing {  
  void Conn.drop() { proceed(); cl  
  void Conn.charge() { ... getTime(  
}
```



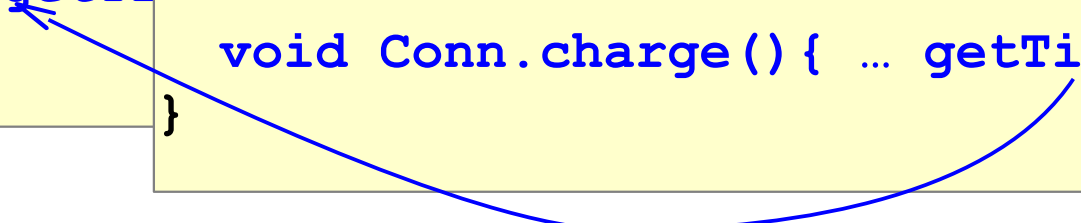
- For **charge ()** in **Billing** to work, baseless partial method **getTime ()** defined in **Timing** should be active beforehand

Telecom example, revisited

```
class Conn {  
    Conn(Custome  
    void complet  
    void drop()  
}
```

```
layer Timing {  
    Timer Conn.tim  
    void Conn.comp  
    void Conn.drop  
    int Conn.getTi  
}
```

```
layer Billing requires Timing {  
    void Conn.drop(){ proceed(); cl  
    void Conn.charge(){ ... getTime(  
}
```



- For **charge ()** in **Billing** to work, baseless partial method **getTime ()** defined in **Timing** should be active beforehand
- In other words, **Billing requires Timing**

Meaning of requires






When layer L **requires** L_1, \dots, L_n

- All of L_1, \dots, L_n (or their sublayers) must have been already activated (in any order) before activating L
- Partial method in L can invoke methods defined in any of L_1, \dots, L_n
- Partial method m in L can proceed when any of L_1, \dots, L_n (or base class) defines m

Coeffect system?

Type Judgment $\Lambda; \Gamma \vdash e : T$

"Under set Λ of activated layers and type env. Γ ,
exp e is given type T "

-  • $\{ \}; c: \text{Conn} \vdash c.\text{getTime}() : \text{int}$
-  • $\{\text{Timing}\}; c: \text{Conn} \vdash c.\text{getTime}() : \text{int}$
-  • $\{ \}; c: \text{Conn} \vdash \text{with } (\text{new Timing}()) c.\text{getTime}() : \text{int}$
-  • $\{ \}; c: \text{Conn} \vdash \text{with } (\text{new Billing}()) c.\text{drop}() : \text{void}$
-  • $\{\text{Timing}\}; c: \text{Conn}$
 $\vdash \text{with } (\text{new Billing}()) c.\text{drop}() : \text{void}$

Inheritance, subtyping and requires

- Sublayer can't require fewer layers than its parent
 - Otherwise, requirement by inherited partial methods may be invalidated
- It seems natural to allow a sublayer to require more layers ...

...Or, maybe not!

```
AbsBilling b =  
    some_condition ? new Billing1() : new Billing2();  
with(b) { ... }
```

- The type system seems to always allow `with(b)` (if `AbsBilling` requires no layer)
- But, what if `Billing2` requires more layers than `AbsBilling`?
 - At run time, dependency is broken!!

Our Solution:

Two subtyping rels for layer types

- Weak subtyping (reflexive transitive closure of extends) for checking `requires` at `with`

```
// L1 extends L2, L3 requires L2  
with(new L1())  
  with(new L3()) { ... }
```

- Normal subtyping (reflexive transitive closure of extends *with invariant requires*) for ordinary subsumption

For more details

- ContextFJ [Hirschfeld, I., Masuhara; FOAL'11]
 - Operational semantics
 - Simple type system disallowing baseless methods
- Type system for baseless methods [I., Hirschfeld, Masuhara; FOOL12]
 - (Slightly different activation semantics)
- Layer inheritance & first-class layers [Inoue&I.; APLAS'15]

Related Work

- Type System for COP [Clarke & Sergey; COP'09]
 - ContextFJ
 - proposed independently of us
 - no inheritance, subtly different semantics
 - Set of method signatures as method-wise dependency information
 - Finer-grained specification
 - No proof of soundness
 - In fact, the type system turns out to be flawed (personal communication), due to `without`

Related Work, contd.

- **Type Systems for Mixins** [Bono et al., Flatt et al., Kamina&Tamai, etc.]
 - Interfaces of classes to be composed
 - Structural type information
 - Composition is fixed once an object is instantiated
 - A similar idea works (to some extent ;-) also for more dynamic composition as in COP
- **Types for FOP, DOP**

Related Work, contd.^2

- Typestate checking [Strom&Yemini'86, etc.]
 - Checking state transition for computational resources (such as files and sockets)
 - Layer configuration can be considered a state

Conclusion

- Dynamic layer composition for describing context-dependent behavioral change concisely and modularly
- Inter-layer dependency (`requires`) works for dynamic composition (as well as static)
- Two kinds of subtyping relations

Future work:

- Type-sound *deactivation*